

Processing and Querying Temporal Data

Johann Gamper

Free University of Bozen-Bolzano, Italy

www.inf.unibz.it/~gamper

Acknowledgments

Joint work with

- M. Böhlen, University of Zurich
- A. Dignös, Free University of Bozen-Bolzano
- C. S. Jensen, Aalborg University

Partially funded by

- Autonomous Province of Bozen-Bolzano
- Swiss National Science Foundation
- EU (Chorochronos)
- unibz (TPG)

- 1 Temporal Databases
- 2 Temporal Alignment – Comprehensive Query Support
- 3 Time Series
- 4 Conclusion

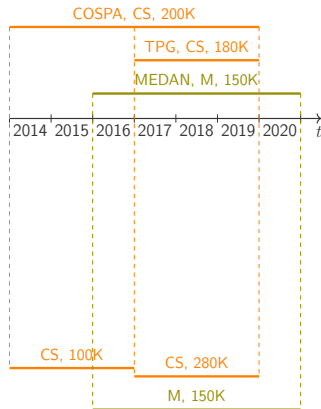
- 1 **Temporal Databases**
- 2 Temporal Alignment – Comprehensive Query Support
- 3 Time Series
- 4 Conclusion

- Project relation with timestamps

Project	Dept	Budget	Time
COSPA	CS	200K	2014 – 2019
TPG	CS	180K	2017 – 2019
MEDAN	M	150K	2016 – 2020

- What is the project budget per department?
 - Independ. of time: \Rightarrow (CS,380K), (M,150K)
 - At the current time (now): \Rightarrow (M,30K)
 - At each time point (sequenced):

Dept	Budget	Time
CS	100K	2014 – 2016
CS	280K	2017 – 2019
M	150K	2016 – 2020



- Timestamps must be **adjusted** for the result
- Some values must be **scaled** to the adjusted timestamps

- Four overlapping phases
 - 1956–1985: **Concept** development
 - 1978–1994: Design of **query languages**
 - 1988–present: **Implementation** aspects
 - Storage and index structures
 - Operator algorithms (join, aggregation)
 - First framework with comprehensive query support
 - 1993–present: **Consolidation** phase
 - Consensus glossary of temporal database concepts
 - Temporal features in SQL
- Still an **active** research area today
 - New application domains need new operations
 - e.g., moving objects, data streams, temporally evolving graphs, etc.

- Period specification for tables (application-time and system-time)

```
CREATE TABLE Emp (  
    Name VARCHAR,  
    Dept VARCHAR,  
    Start DATE,  
    End DATE,  
    PERIOD FOR Period (Start, End) );
```

- Temporal UPDATE/DELETE behavior
- Temporal primary and foreign keys

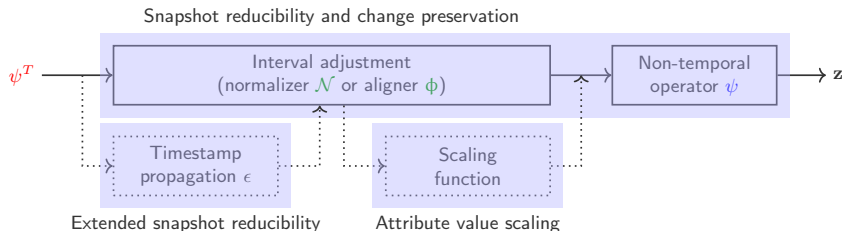
```
ALTER TABLE Emp  
ADD PRIMARY KEY (Name, Period WITHOUT OVERLAPS)
```

- Predicates/Functions for periods to support querying

- Support for storage and update
- Limited support for query formulation!

- 1 Temporal Databases
- 2 Temporal Alignment – Comprehensive Query Support**
- 3 Time Series
- 4 Conclusion

- Reduce temporal operators ψ^T to nontemporal operators ψ
 - Adjust time periods of input relations
 - Apply non-temporal operator



- Minimal changes to DBMS: normalizer and aligner primitives
- Existing query optimization/indexing works

- What is the budget per department: $D^{\vartheta T}SUM(Budget)(\mathbf{P})$

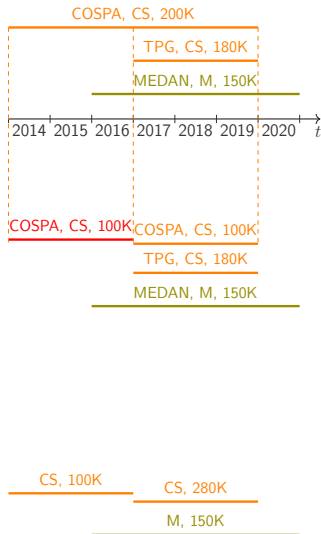
Project	Dept	Budget	Time
COSPA	CS	200K	2014 – 2019
TPG	CS	180K	2017 – 2019
MEDAN	M	150K	2016 – 2020

adjustment & scaling

Project	Dept	Budget	Time
COSPA	CS	100K	2014 – 2016
COSPA	CS	100K	2017 – 2019
TPG	CS	180K	2017 – 2019
MEDAN	M	150K	2016 – 2020

nontemp. aggregation by Dept and Time

Dept	Sum	Time
CS	100K	2014 – 2016
CS	280K	2017 – 2019
M	150K	2016 – 2020



Reduction Rules: $\psi^T \longrightarrow \{\mathcal{N}, \phi\} + \psi$

Operator	Reduction
Selection	$\sigma_{\theta}^T(\mathbf{r}) = \sigma_{\theta}(\mathbf{r})$
Projection	$\pi_{\mathbf{B}}^T(\mathbf{r}) = \pi_{\mathbf{B}, \mathbf{T}}(\mathcal{N}_{\mathbf{B}}(\mathbf{r}, \mathbf{r}))$
Aggregation	$\mathbf{B} \vartheta_F^T(\mathbf{r}) = \mathbf{B}, \mathbf{T} \vartheta_F(\mathcal{N}_{\mathbf{B}}(\mathbf{r}, \mathbf{r}))$
Difference	$\mathbf{r} -^T \mathbf{s} = \mathcal{N}_{\mathbf{A}}(\mathbf{r}, \mathbf{s}) - \mathcal{N}_{\mathbf{A}}(\mathbf{s}, \mathbf{r})$
Union	$\mathbf{r} \cup^T \mathbf{s} = \mathcal{N}_{\mathbf{A}}(\mathbf{r}, \mathbf{s}) \cup \mathcal{N}_{\mathbf{A}}(\mathbf{s}, \mathbf{r})$
Intersection	$\mathbf{r} \cap^T \mathbf{s} = \mathcal{N}_{\mathbf{A}}(\mathbf{r}, \mathbf{s}) \cap \mathcal{N}_{\mathbf{A}}(\mathbf{s}, \mathbf{r})$
Cart. Prod.	$\mathbf{r} \times^T \mathbf{s} = \alpha(\phi_{\mathbf{T}}(\mathbf{r}, \mathbf{s}) \bowtie_{\mathbf{r}, \mathbf{T}=\mathbf{s}, \mathbf{T}} \phi_{\mathbf{T}}(\mathbf{s}, \mathbf{r}))$
Inner Join	$\mathbf{r} \bowtie_{\theta}^T \mathbf{s} = \alpha(\phi_{\theta}(\mathbf{r}, \mathbf{s}) \bowtie_{\theta \wedge \mathbf{r}, \mathbf{T}=\mathbf{s}, \mathbf{T}} \phi_{\theta}(\mathbf{s}, \mathbf{r}))$
Left O. Join	$\mathbf{r} \Join_{\theta}^T \mathbf{s} = \alpha(\phi_{\theta}(\mathbf{r}, \mathbf{s}) \Join_{\theta \wedge \mathbf{r}, \mathbf{T}=\mathbf{s}, \mathbf{T}} \phi_{\theta}(\mathbf{s}, \mathbf{r}))$
Right O. Join	$\mathbf{r} \Join_{\theta}^T \mathbf{s} = \alpha(\phi_{\theta}(\mathbf{r}, \mathbf{s}) \Join_{\theta \wedge \mathbf{r}, \mathbf{T}=\mathbf{s}, \mathbf{T}} \phi_{\theta}(\mathbf{s}, \mathbf{r}))$
Full O. Join	$\mathbf{r} \Join_{\theta}^T \mathbf{s} = \alpha(\phi_{\theta}(\mathbf{r}, \mathbf{s}) \Join_{\theta \wedge \mathbf{r}, \mathbf{T}=\mathbf{s}, \mathbf{T}} \phi_{\theta}(\mathbf{s}, \mathbf{r}))$
Anti Join	$\mathbf{r} \triangleright_{\theta}^T \mathbf{s} = \phi_{\theta}(\mathbf{r}, \mathbf{s}) \triangleright_{\theta \wedge \mathbf{r}, \mathbf{T}=\mathbf{s}, \mathbf{T}} \phi_{\theta}(\mathbf{s}, \mathbf{r})$

Temporal Op.

= Primitive + Traditional Op.

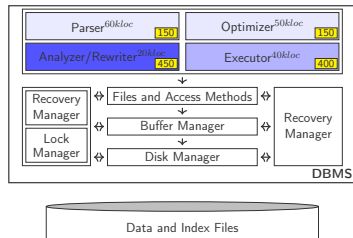
- Integration in DBMS kernel (approx. 1000 LOC)
- Dignös et al.: *Extending the kernel of a relational DBMS with comprehensive support for sequenced temporal queries*. TODS, 2016.



- Submitted as patch to PostgreSQL

- For more information go to: <http://tpg.inf.unibz.it>

- Ongoing project: we seek PhD students and PostDocs to [join!](#)



- Boosting performance
 - Customized alignment primitives to reduce intermediate relation size
 - Precise cost estimates based on temporal distribution
 - Equivalence rules for the interaction of primitives with RA
- Support for multiple time dimensions
 - e.g., valid time and transaction time (supported by SQL:2011)
- SQL extension to facilitate the formulation of temporal queries

- 1 Temporal Databases
- 2 Temporal Alignment – Comprehensive Query Support
- 3 Time Series**
- 4 Conclusion

- A **special type** of temporal data: **sequence of point values**
- Analysis has to consider **entire** time series (not just single point)
- Mostly based on **similarity**: find most similar TS is a fundamental query

Databases

Operations

Selection Projection Aggregation
Join Difference Intersection Union

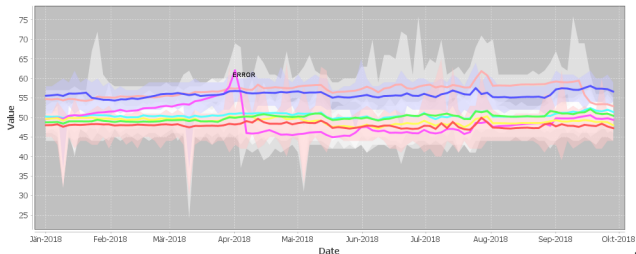
Time series

Alignment Missing value imputation
Outlier detection Resampling Dimensionality reduction
Aggregation Anomaly detection Classification Compression
Correlation analysis Discord detection Forecasting
Function approximation Motif discovery Prediction
Predictive maintenance Seasonality analysis Segmentation
Similarity search Subsequence search etc.

Technologies

RA
SQL
RDBMS
B-tree
Hash
ED
DTW
PAA
PLA
DFT
iSAX
ADS
SQL-TS
TSMS
⋮

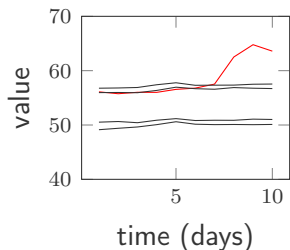
- **Goal:** Predict device errors or maintenance steps
- **Idea:** Spot patterns in sensor data with high prognostic accuracy



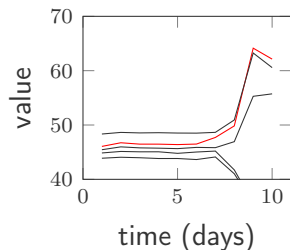
Preprocessing

- is essential for accurate/reliable analytics
- is work-intensive, ad-hoc, lack of methods to steer the process

Leads to an error

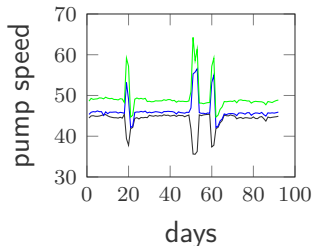


Does not lead to an error

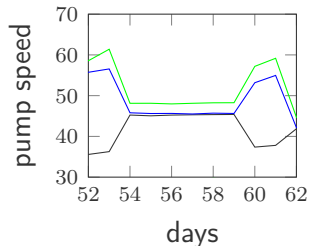


- Error pattern becomes only visible by looking at **multivariate signals**
- How to determine a **good subset of signals**?

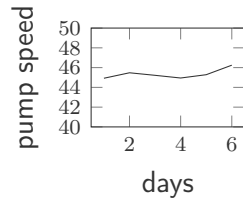
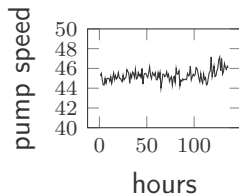
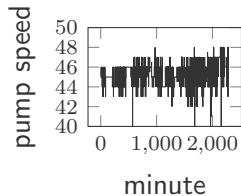
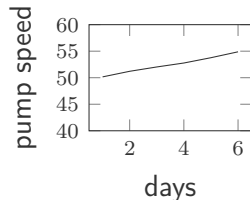
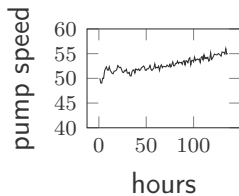
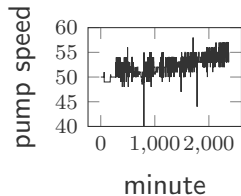
Signals with three errors



Window covers two errors



- Different **window lengths** might or might not detect an error
- How to determine a **good window length**?



- Different **granularities/representations** emphasize different information
- How to determine a **good granularity/representation**?

- Investigation of **systematic preprocessing** techniques and methods
 - Impact on downstream analysis
- **Time series management system (TSMS)**
 - Comprehensive support for time series management/analytics/mining
 - Integrate into RDBMs?

- 1 Temporal Databases
- 2 Temporal Alignment – Comprehensive Query Support
- 3 Time Series
- 4 Conclusion**

Temporal databases

- Temporal features in [SQL:2011](#)
- [Query support](#) is largely missing
- [Temporal alignment](#) offers comprehensive query support

Time series

- [Special kind of temporal data](#), heavily based on [similarity](#)
- [Operations](#) studied in isolation
- [Preprocessing](#) not well studied

Future work

- [Performance/query optimization](#)
- [Multiple](#) time dimensions
- [SQL extension](#)

Future work

- Systematic investigation of [preprocessing](#)
- [\(Relational\) Time series management system](#)

Thank You!