

Identification of a 2-additive bi-capacity by using mathematical programming

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- In many situations, human decision makings are easier using bipolar scales (negative, neutral and positive parts) :
 - ▶ Bi-capacities (BC) allows extending CI to bipolar scales
 - ▶ But the nb of weights to set is bigger : $3^n - 1$

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- We propose optimization problems to elicit the values of a 2A-BC
- Our framework presents several common points with supervised machine learning tasks and we underline these points

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- 1 Bi-capacities and bipolar Choquet integrals
- 2 Identifying a 2-additive bi-capacity
- 3 An illustrative example
- 4 Relationships with supervised learning tasks in machine learning
- 5 Conclusion and future work

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- On 3^N , the relation \sqsubseteq is such that $(A_1, A_2) \sqsubseteq (B_1, B_2) \Leftrightarrow A_1 \subseteq B_1 \wedge B_2 \subseteq A_2$

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- 3^N is given by pairs indicated by \checkmark in the following indicator matrix :

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- A set function $\nu : 3^N \rightarrow \mathbb{R}$ is a **bi-capacity (BC)** on 3^N if it satisfies the following two conditions [Grabisch and Labreuche, 2005b], [Grabisch and Labreuche, 2008] :

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- ν is said to be normalized if in addition, it holds :

$$\nu(N, \emptyset) = 1 \wedge \nu(\emptyset, N) = -1$$

Bipolar Möbius transform (BMT) of a BC

- A BC ν can be associated to its **bipolar Möbius transform (BMT)** denoted b and defined by [Fujimoto, 2004, Fujimoto, 2007] :

$$\begin{aligned}
 b(A_1, A_2) &:= \sum_{\substack{B_1 \subseteq A_1 \\ B_2 \subseteq A_2}} (-1)^{|A_1 \setminus B_1| + |A_2 \setminus B_2|} \nu(B_1, B_2) \\
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- There is a one-to-one relation between ν and b and the converse relation is given by :

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B1 Note that the property $\nu(\emptyset, \emptyset) = 0$ translates into :

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 - $b = 0$ if the nb of criteria involved in (A_1, A_2) is greater than k :

$$\forall (A_1, A_2) \in 3^N : |A_1 \cup A_2| > k \Rightarrow b(A_1, A_2) = 0$$

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B3 $b \neq 0$ for at least one pair such that the nb of criteria involved is k :

$$\exists (A_1, A_2) \in 3^N : |A_1 \cup A_2| = k \wedge b(A_1, A_2) \neq 0$$

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- In this case, only subsets with at most two criteria matter
- To lighten the notations we will take :

$$b_{i|} = b(\{i\}, \emptyset) \quad ; \quad b_{ij|} = b(\{i, j\}, \emptyset) \quad ; \quad b_{i|j} = b(\{i\}, \{j\})$$

Example of the BMT of a 2A-BC

- $N = \{1, 2, 3\}$
- We assume b is a BMT of a 2A-BC ν
- \checkmark are (possibly) non null elements and \times are (necessarily) null elements :

$$b = \begin{array}{c} (A, B) \\ \emptyset \\ 1 \\ 2 \\ 3 \\ 12 \\ 13 \\ 23 \\ 123 \end{array} \begin{pmatrix} \checkmark & \checkmark & \checkmark & \checkmark & \checkmark & \checkmark & \checkmark & \times \\ \checkmark & \cdot & \checkmark & \checkmark & \cdot & \cdot & \times & \cdot \\ \checkmark & \checkmark & \cdot & \checkmark & \cdot & \times & \cdot & \cdot \\ \checkmark & \checkmark & \checkmark & \cdot & \times & \cdot & \cdot & \cdot \\ \checkmark & \cdot & \cdot & \times & \cdot & \cdot & \cdot & \cdot \\ \checkmark & \cdot & \times & \cdot & \cdot & \cdot & \cdot & \cdot \\ \checkmark & \times & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \times & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \end{pmatrix}$$

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- The nb of elements to be set for b reduces from 3^n to $2n^2 + 1$ (27 vs 19 in this example)

2-additive BC (cont'd)

P1 The **monotonicity** property reduces to [Mayag et al, 2012] :

$$\forall (A, B) \in 3^N, \forall k \in A : b_{k|} + \sum_{j \in B} b_{k|j} + \sum_{i \in A \setminus k} b_{i|k} \geq 0$$

$$\forall (A, B) \in 3^N, \forall k \in A : b_{|k} + \sum_{j \in B} b_{j|k} + \sum_{i \in A \setminus k} b_{|ik} \leq 0$$

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P2 The **normalization** property reduces to [Mayag et al, 2012] :

$$\sum_{i \in N} b_{i|} + \sum_{\{i,j\} \subseteq N} b_{ij|} = 1$$

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Bipolar Choquet integral (BCI) wrt a 2A-BC

P3 The **bipolar Choquet integral (BCI)** wrt b denoted \mathcal{C}_b is given by :

$$\begin{aligned} \mathcal{C}_b(x) = & \sum_{i=1}^n b_{|i|} x_i^+ + \sum_{i=1}^n b_{|i|} x_i^- + \sum_{i,j=1}^n b_{|ij|} (x_i^+ \wedge x_j^-) \\ & + \sum_{\{i,j\} \subseteq N} b_{ij|} (x_i^+ \wedge x_j^+) + \sum_{\{i,j\} \subseteq N} b_{|ij|} (x_i^- \wedge x_j^-) \end{aligned}$$

$$\text{where } x_i^+ = \begin{cases} x_i & \text{if } x_i > 0 \\ 0 & \text{if } x_i \leq 0 \end{cases} \quad \text{and } x_i^- = \begin{cases} -x_i & \text{if } x_i < 0 \\ 0 & \text{if } x_i \geq 0 \end{cases}$$

$$\text{and } a \wedge b = \min(a, b)$$

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 - ▶ the scores distribution $x = (x_1, \dots, x_n)$
 - ▶ the overall score $S(x)$ that is the aggregated value for x

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- The preferences model of the DM is modeled by a 2A-BC ν represented by its BMT b
- The DM provides a subset of alternatives $X' \subseteq X$ and $\forall x \in X'$ we are given :
 - ▶ the scores distribution $x = (x_1, \dots, x_n)$
 - ▶ the overall score $S(x)$ that is the aggregated value for x
- Let us denote by T the set of pairs $\{((x_1, \dots, x_n), S(x))\}_{x \in X'}$

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- No extra information is provided
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Taking into account preference relations

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- However it might happen that it is impossible to satisfy these constraints for all pairs :
 - ▶ Our restriction to 2A-BC might be too strong and more general BC could better fit the problem
 - ▶ The DM could provide scores and preferences that present incoherences and which are not fully representable by any BC

Taking into account preference relations (cont'd)

C2 If C1 cannot be satisfied, we replace it with more flexible constraints :

$$\forall x, x' \in X', x \neq x' : S(x) - S(x') \geq 0 \Rightarrow C_b(x) - C_b(x') \geq -\xi_{xx'}$$

where $\xi_{xx'} \geq 0$ are **slack variables** of the model

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L1 We could also impose C_b to be bounded :

$$\forall x \in X' : lb \leq C_b(x) \leq ub$$

The split approach

C3 We add a variable $\varepsilon \geq 0$ that reflects the difference between $\mathcal{C}_b(x)$ and $\mathcal{C}_b(x')$:

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- This is a linear program

A flexible version of the split approach

- In case of inconsistencies, we use the following objective function :

$$\max \left(\varepsilon - \sum_{x, x': S(x) \geq S(x')} \xi_{xx'} \right)$$

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- ▶ C4 : preference relations “with $\varepsilon - \xi_{xx'}$ ”

$$\forall x, x' \in X', x \neq x' : S(x) - S(x') \geq 0 \Rightarrow \mathcal{C}_b(x) - \mathcal{C}_b(x') \geq \varepsilon - \xi_{xx'}$$

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- We could also minimize the residual sum of squares
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 - ▶ L1 : C_b bounded
 - ▶ C1 : preference relations
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- In case of incoherences, we optimize the following objective function :

$$\min \left(\sum_{x \in X'} (S(x) - C_b(x))^2 + \sum_{x, x': S(x) \geq S(x')} \xi_{xx'} \right)$$

- subject to :

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An example without incoherence

- It concerns the grades (scores) obtained by 7 students (alternatives) for $n = 5$ subjects (criteria) : statistics (S), probability (P), economics (E), management (M), and English (En).
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Stu.	S	P	E	M	En
<i>a</i>	4	-3	-3	-3	4
<i>b</i>	4	-3	4	-3	-3
<i>c</i>	-3	-3	4	-3	4
<i>d</i>	4	4	-3	-3	-3
<i>e</i>	-3	-3	4	4	-3
<i>f</i>	-3	-3	4	-3	-3
<i>g</i>	-3	-3	-3	-3	4

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Stu.	S	P	E	M	En	S
<i>a</i>	4	-3	-3	-3	4	1
<i>b</i>	4	-3	4	-3	-3	0.5
<i>c</i>	-3	-3	4	-3	4	0
<i>d</i>	4	4	-3	-3	-3	-0.5
<i>e</i>	-3	-3	4	4	-3	-1
<i>f</i>	-3	-3	4	-3	-3	-1.5
<i>g</i>	-3	-3	-3	-3	4	-2

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Stu.	S	P	E	M	En	S	<i>split</i>
<i>a</i>	4	-3	-3	-3	4	1	1.68
<i>b</i>	4	-3	4	-3	-3	0.5	1.04
<i>c</i>	-3	-3	4	-3	4	0	0.41
<i>d</i>	4	4	-3	-3	-3	-0.5	-0.23
<i>e</i>	-3	-3	4	4	-3	-1	-0.86
<i>f</i>	-3	-3	4	-3	-3	-1.5	-1.5
<i>g</i>	-3	-3	-3	-3	4	-2	-2.14

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Stu.	S	P	E	M	En	S	$split$	$split$ $flex$
<i>a</i>	4	-3	-3	-3	4	1	1.68	1.02
<i>b</i>	4	-3	4	-3	-3	0.5	1.04	0.38
<i>c</i>	-3	-3	4	-3	4	0	0.41	-0.25
<i>d</i>	4	4	-3	-3	-3	-0.5	-0.23	-0.89
<i>e</i>	-3	-3	4	4	-3	-1	-0.86	-1.53
<i>f</i>	-3	-3	4	-3	-3	-1.5	-1.5	-2.16
<i>g</i>	-3	-3	-3	-3	4	-2	-2.14	-2.8

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Stu.	S	P	E	M	En	S	$split$	$split$ $flex$	rss
<i>a</i>	4	-3	-3	-3	4	1	1.68	1.02	1
<i>b</i>	4	-3	4	-3	-3	0.5	1.04	0.38	0.5
<i>c</i>	-3	-3	4	-3	4	0	0.41	-0.25	0
<i>d</i>	4	4	-3	-3	-3	-0.5	-0.23	-0.89	-0.5
<i>e</i>	-3	-3	4	4	-3	-1	-0.86	-1.53	-1
<i>f</i>	-3	-3	4	-3	-3	-1.5	-1.5	-2.16	-1.5
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<i>g</i>	-3	-3	-3	-3	4	-2	-2.14	-2.8	-2	-2

An example with an incoherence

- We change S into S' and the only difference is $S'(g) = 0.5$ while $S(g) = -2$. There is an incoherence between c and g : $S'(g) > S'(c)$ while g is Pareto dominated by c

Stu.	S	P	E	M	En	S'
<i>a</i>	4	-3	-3	-3	4	1
<i>b</i>	4	-3	4	-3	-3	0.5
<i>c</i>	-3	-3	4	-3	4	0
<i>d</i>	4	4	-3	-3	-3	-0.5
<i>e</i>	-3	-3	4	4	-3	-1
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<i>g</i>	-3	-3	-3	-3	4	0.5

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<i>c</i>	-3	-3	4	-3	4	0	.
<i>d</i>	4	4	-3	-3	-3	-0.5	.
<i>e</i>	-3	-3	4	4	-3	-1	.
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<i>g</i>	-3	-3	-3	-3	4	0.5	.

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<i>a</i>	4	-3	-3	-3	4	1	.	0.22
<i>b</i>	4	-3	4	-3	-3	0.5	.	-0.28
<i>c</i>	-3	-3	4	-3	4	0	.	-0.78
<i>d</i>	4	4	-3	-3	-3	-0.5	.	-1.28
<i>e</i>	-3	-3	4	4	-3	-1	.	-1.78
<i>f</i>	-3	-3	4	-3	-3	-1.5	.	-2.28
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<i>b</i>	4	-3	4	-3	-3	0.5	.	-0.28	.	0.62
<i>c</i>	-3	-3	4	-3	4	0	.	-0.78	.	0.12
<i>d</i>	4	4	-3	-3	-3	-0.5	.	-1.28	.	-0.5
<i>e</i>	-3	-3	4	4	-3	-1	.	-1.78	.	-1
<i>f</i>	-3	-3	4	-3	-3	-1.5	.	-2.28	.	-1.5
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<i>d</i>	4	4	-3	-3	-3	-0.5	.	-1.28	.	-0.5
<i>e</i>	-3	-3	4	4	-3	-1	.	-1.78	.	-1
<i>f</i>	-3	-3	4	-3	-3	-1.5	.	-2.28	.	-1.5
<i>g</i>	-3	-3	-3	-3	4	0.5	.	-0.78	.	0.12

- For *split flex* and *rss flex*, $\xi_{gc} > 0$ while for other pairs the slack variables are null

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Relationships with supervised Learning (SL)

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- The set T is similar as a training set in SL
- Eliciting the parameters of the preference model (b) and the aggregation operator (\mathcal{C}_b) is similar as inferring the parameters of a SL model (like the coefficients of the linear regression)

Relationships with supervised Learning (SL)

The addressed task has many similarities with supervised learning (SL) tasks in machine learning (ML)

- The set T is similar as a training set in SL
- Eliciting the parameters of the preference model (b) and the aggregation operator (C_b) is similar as inferring the parameters of a SL model (like the coefficients of the linear regression)
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- 2A-BC is a less vast family of preference models than unconstrained BC. Thus constraining the BC to be 2-additive is like choosing a family of hypothesis with potentially greater bias but lower variance (we expect better generalization for unseen alternatives)

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$$\forall (A_1, A_2) \in 3^N : |A_1 \cup A_2| > k \Rightarrow b(A_1, A_2) = 0$$

B3 $b \neq 0$ for at least one pair such that the nb of criteria involved is k :

$$\exists (A_1, A_2) \in 3^N : |A_1 \cup A_2| = k \wedge b(A_1, A_2) \neq 0$$

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- Thus, the model can be a 1-additive BC and not necessarily a 2A-BC
- But, this is not guaranteed : we need to add a penalty term to favor sparse solutions
- The slack variables $\xi_{xx'}$ and the constraints C1, C2, C3, C4 have relationships with SVM (margin, structured SVM)

Outline

- 1 Bi-capacities and bipolar Choquet integrals
- 2 Identifying a 2-additive bi-capacity
- 3 An illustrative example
- 4 Relationships with supervised learning tasks in machine learning
- 5 Conclusion and future work

Conclusion and future work







- We proposed to use the BCI wrt a 2A-BC as a preference model in MCDM
- We introduced two kinds of optimization problems to elicit a 2A-BC
- Our models allows dealing with inconsistencies
- Our setting and elicitation model has several common points with supervised learning

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- We introduced two kinds of optimization problems to elicit a 2A-BC
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- As for ongoing and future work :
 - ▶ Further exploit ML concepts in MCDM like adding a penalty term to have sparse b
 - ▶ Have a better understanding of the behavior of b and ν provided by the different methods
 - ▶ Extend our elicitation framework to integrate other information provided by the DM such as the importance or the interaction between criteria

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Thank you for your attention ! Questions ?